

- Suppose an array is sorted,
and you'd like to find an entry in that array
- You could use a linear search from a previous topic
- A linear search is necessary if an array is not sorted, for the entry you're looking for may be anywhere
- Question: can we speed up the search if the array is sorted?
- Imagine if you had a book, and you had to find page 147
- If the book was 308 pages, would you start with page 1 ?
- Like our linear search,
the binary search will have a similar function declaration
std::size_t binary_search( double const array[],
std::size_t const capacity double const value );


##  <br> Implementing the binary search

- We will start by checking the array is sorted,
and setting the bounds
assert( is_sorted( array, capacity ) == capacity );

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##  <br> Searching sorted arrays

- Suppose you are searching an array with capacity 100 for 9237.1 - The value could be anywhere between index 0 and 99
- Suppose you guess the index $(0+99) / 2==49$
- Check if array[49] == value
- In the off chance this entry equals the value, great!
- Check if array[49] < 9237.1
- We're too low, so we are restricted to 50 to 99
- Next time, inspect index $(50+99) / 2==74$
- Otherwise, array[49] > 9237.1
- We're too high, so we are restricted to 0 to 48
- Next time, inspect index $(0+48) / 2==24$
- We cannot use a for loop, so we'll have to use a while loop
- We'll hold off on the condition
assert( is_sorted( array, capacity ) == capacity );
std::size_t lower_index\{0\};
std::size_t upper_index\{capacity - 1\};



##  <br> Implementing the binary search

- First, we will calculate the average of the lower and upper indices

```
assert( is_sorted( array, capacity ) == capacity );
std::size_t lower_index{0};
std::size_t upper_index{capacity - 1}
while (... ) {
    std::size_t average_index{ (lower_index + upper_index)/2 };
```

    \}
    - Next, we will check if this entry contains what we are searching for
assert( is_sorted( array, capacity ) == capacity );
std::size_t lower_index\{0\};
std::size_t upper_index\{capacity - 1$\}$;
while (... ) \{
std::size_t average_index\{ (lower_index + upper_index)/2 \};
f ( array[average_index] == value ) \{ return average_index;
\}
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##  <br> Implementing the binary search

- Otherwise, the only other possibility is that the entry is too large - We'll update the upper index
assert( is_sorted( array, capacity ) == capacity );
std::size_t lower_index\{0\};
std::size_t upper_index\{capacity - 1 \};
while ( ... ) \{
std::size_t average_index\{ (lower_index + upper_index)/2 \};
if ( array[average_index] == value ) \{ return average_index;
\} else if ( array[average_index] < value ) \{
lower_index $=$ average_index +1 :
upper_index = average_index - 1 ;

```
}
```

- Now for the more interesting question:
- When do we stop searching?
assert( is_sorted( array, capacity ) $==$ capacity );
std::size_t lower_index\{0\};
std::size_t upper_index\{capacity - 1$\}$;
while $(\ldots)$ \{
td: : size t average index\{ (lower index + upper index) $/ 2$
if ( array[average_index] == value ) \{
return average_index;
\} else if ( array[average_index] < value ) \{
lower_index = average_index +1 ;
\} else \{
upper_index = average_index - 1 ;
\}


##  <br> While loop condition

- Question: What is the necessary condition for the while loop?
- If we are searching for something in the array,
no halting condition is necessary
- This algorithm is guaranteed to find the entry
- Thus, the condition would be
while ( true ) \{
// Find the average and update accordingly
- Question: What is the necessary halting condition?
- What if we are searching for something not in the array?
- With every iteration of the loop where we don't find the value, either
- The value of lower_index will be increasing, or
- The value of upper_index will be decreasing


## 囷 <br> While loop condition

- Question: What is the necessary halting condition?
- What if we are searching for something not in the array?


## Case 1

- Suppose that at one step, lower_index == upper_index
- If this is the case, average_index will equal this value, so if array[average_index] < value,
uppex_index == lower_index - 1
and if array [average_index] > value,
lower_index == upper_index + 1
- In either case, lower_index > upper_index
- Question: What is the necessary halting condition?
- What if we are searching for something not in the array?


## Case 2

- Suppose that at one step, lower_index + 1 == upper_index
- If this is the case, average_index will equal lower_index, so if array[average_index] < value,
uppex_index == lower_index - 1
and if array[average_index] > value,
lower_index == upper_index
- In the first case, lower_index > upper_index and in the second, we are back to Case 1
- Thus, we should continue iterating so long as:
while (lower_index <= upper_index) \{
// Find the average and update accordingly
\}
return capacity;
- Thus, we should stop looping as soon as
lower_index > upper_index


## While loop condition



## While loop condition

- Question: What is the necessary halting condition?
- What if we are searching for something not in the array?


## Case 3

- Suppose that at one step, lower_index $+2==$ upper_index
- If this is the case, average_index will equal lower_index + 1, so if array[average_index] < value,
and if array[average_index] > value,
lower_index == upper_index
- In both possibilities, we are back to Case 1


- Here is a test:
std::size_t const CAPACITY\{5\};
double data[CAPACITY]\{ $0.0,1.2,1.5,2.7,4.6\}$

```
(dont k{0}; k <= CAPACITY*10; ++k )
double x{ k/10.0 };
std::size_t index{ binary_search( data, CAPACITY, x ) }
assert( (0 <= index) && (index <= CAPACITY));
if ( index == CAPACITY ) {
            std::cout << x << " not found" << std::endl
        } else {
            std::cout << x << "\t== data[" << index << "] = 
            std::cout << x << "\t== data[" << index
```

                    [
    \} return $\theta$;
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##  <br> Weakness in our test

- Did any of you notice a weakness in the test?
int main() \{
std::size t const CAPACITY\{5\};
double data[CAPACITY]\{ 0.0, 1.2, 1.5, 2.7, 4.6 \};
for ( int k\{0\}; k <= CAPACITY*10; ++k ) double $\times\{k / 10.0$ \};

We never searched for a value less than the first entry

- Here is the output:


- We can fix this:
int main() \{
double data[CAPACITY]\{ 0.3, 1.2, 1.5, 2.7, 4.6 \};
for ( int k\{0\}; k <= CAPACITY*10; ++k ) \{ double $x\{k / 10.0\}$;
- This should indicate 0,0.1 and 0.2 are not found
- On my computer, however, I get:

Segmentation fault (core dumped)

##  <br> Weakness in our implementation

- If both lower_index and upper_index equal 0,
average_index = 0,
and if array[0] > value, then
upper_index = 0 - 1;
- Problem: std::size_t is unsigned, so 0 - 1 causes a carry, so now upper_index == 0xffff...f
- How do we test for this?
- Note that this maximum value is greater than or equal to the capacity of the array
- Thus, if upper_index == 0xffffl..f
then upper_index >= capacity

Our complete implementation
std::size_t binary_search( double const array[], std::size_t const capacity double const value ) \{
assert( is_sorted( array, capacity ) == capacity );
std:: size_t lower_index\{ 0$\}$;
std::size_t upper_index\{capacity - 1\}
while ( (lower_index <= upper_index)
\&\& (upper_index < capacity)) \{
std::size_t average_index\{ (lower_index + upper_index)/2 \};
if ( array[average_index] == value ) \{ return average_index;
\} else if (array[average_index] < value ) \{
lower_index = average_index + 1;
\} else \{
upper_index $=$ average_index - 1 ;
\}

##  <br> Weakness in our implementation

- Thus we should not continue if either of these conditions is true: lower_index > upper_index upper_index >= capacity
- Thus we should continue if both of these conditions are false: lower_index > upper_index upper index >= capacity
- Thus we should continue if both of these conditions are true: lower_index <= upper_index upper_index < capacity

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##  <br> What type of error is this?

- This error in our function was not a problem with binary search
- If upper index could take on the value -1,
the implementation would work as expected
- It is because the index is unsigned that this error appears
- We call such an error a semantic error
- We expect integer arithmetic to work as it does in the real world
- This is not what happens with unsigned integers in C++
- Interesting observation:
- This fix also fixes the problem if capacity is 0
- The initial value of upper_limit is 0 - 1
- We seem to have taken care of two weaknesses with one fix

How often must we search the array?

- Question:

What is the maximum number of entries of the array that we must inspect?

- Our test requires that we inspect the output
- Can we write a test that does not require us to do so? int main() \{
std:: :size_t const CAPACITY\{5\};
double datal CAPACITY1 std: :size_t found_count $\{0\}$;
for (int $k\{\theta\} ; k<=$ CAPACITY*10; ++k ) \{
- 

std::size $t$ indexf binary_search( data, CAPACITY, x ) \};
assert( ( $0<=$ index) \&\& (index <= CAPACITY) );
if (index $!=$ CAPACITY) $\left\{\begin{array}{l}\text { assert }(\operatorname{data}[\text { index }]=x\end{array}\right.$
++ found_count;
\}



- Suppose the array has a capacity of 127
- In this case, we must search entries from 0 to 126
- We are searching 127 entries
- If it is not at index 63
- We must search the indices from 0 to 62 or 64 to 126
- In both cases, we are restricted to searching 63 entries
- In the first case, if it is not at index 31
- We must search the indices from 0 to 30 or 32 to 62
- In the second case, if it is not at index 95
- We must search the indices from 64 to 94 or 96 to 126
- In all four cases, we are restricted to searching 31 entries
- Note these values are $2^{7}-1=127,2^{6}-1=63,2^{5}-1=31$,
- You may correctly deduce that this pattern will continue:

$$
2^{4}-1=15, \quad 2^{3}-1=7, \quad 2^{2}-1=3,2^{1}-1=1
$$

- In your course on algorithms and data structures,
you will prove that a binary search will insepct no more than $\log _{2}(n)+1$
entries of the array
- A linear search on an array of capacity one million may require up to searching one million entries
- A binary search on a sorted array of capacity one million will require no more than $\log _{2}(1000000)+1=20.93156857$
- That is, inspecting no more than 20 entries of the array

[1] Wikipedia,
https://en.wikipedia.org/wiki/Binary_search
[2] Dictionary of Algorithms and Data Structures (DADS)
https://xlinux.nist.gov/dads/HTML/binarySearch.html


##  <br> 

- Following this presentation, you now:
- Understand how to implement a binary search
- Are aware that you must thoroughly test your implementation
- Understand that there may be issues with implementing the algorithm as described
- In our case, for an unsigned integer, we calculated $\theta-1$,
which resulted not in -1 , but rather at 0xfff...f
- Are aware that a binary search is relatively fast compared to a linear search, but the array must be sorted


Proof read by Dr. Thomas McConkey and Charlie Liu.

These slides were prepared using the Georgia typeface. Mathematical equations use Times New Roman, and source code is presented using Consolas.

The photographs of lilacs in bloom appearing on the title slide and accenting the top of each other slide were taken at the Royal Botanical Gardens on May 27, 2018 by Douglas Wilhelm Harder. Please see
https://www.rbg.ca/


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[^0]:    std: :size_t lower_index\{0\};
    std::size_t upper_index\{capacity - 1\}

